

DNS Inconsistency

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Introduction

- Parent NS RRset ***p***

```
example 1D IN NS ns1.example.edu.  
example 1D IN NS ns2.example.edu.
```

- Child NS RCset ***c***

```
example 1D IN NS ns1.example.edu.  
example 1D IN NS ns2.example.edu.  
example 1D IN NS ns3.example.edu.
```

Real, Badly Inconsistent Example

- Child NS RRset *cdm.depaul.edu*

```
cdm.depaul.edu. 3600 IN NS ns1.cti.depaul.edu.  
cdm.depaul.edu. 0 IN NS shemp.cti.depaul.edu.  
cdm.depaul.edu. 3600 IN NS ns-colo.cti.deapaul.edu.  
cdm.depaul.edu. 3600 IN NS dc-colo-cti.cti.depaul.edu.  
cdm.depaul.edu. 3600 IN NS bach.cti.depaul.edu.  
cdm.depaul.edu. 3600 IN NS ellington.cti.depaul.edu.  
cdm.depaul.edu. 3600 IN NS moe.cti.depaul.edu.  
cdm.depaul.edu. 3600 IN NS mozart.cti.depaul.edu.  
  
. . .  
ns-colo.cti.depaul.edu. AAAA 2002:d8dc:b452::dbdc:b452  
dc-colo-cti.cti.depaul.edu. A 10.128.30.2
```

Conjecture

DNS infrastructure (NS parent/child RRset) inconsistency arises from **asynchronous** and **uncoordinated** NS RRset **configuration**

Methodology: .edu traversal

- Obtain all .edu names using `whois` * hack

```
for each name in edu
  mark root_servers as visited
  get NS_RRset from an .edu NS for the name
  for each S in NS_RRset
    do_query( name, S )
```

`do_query:`

```
  return if already queried S for name
  mark ( name, S ) as visited
  get new_NS_RRset for name from S
  for each S* in new_NS_RRset
    do_query( name, S* )
```

NS name mapping ambiguity

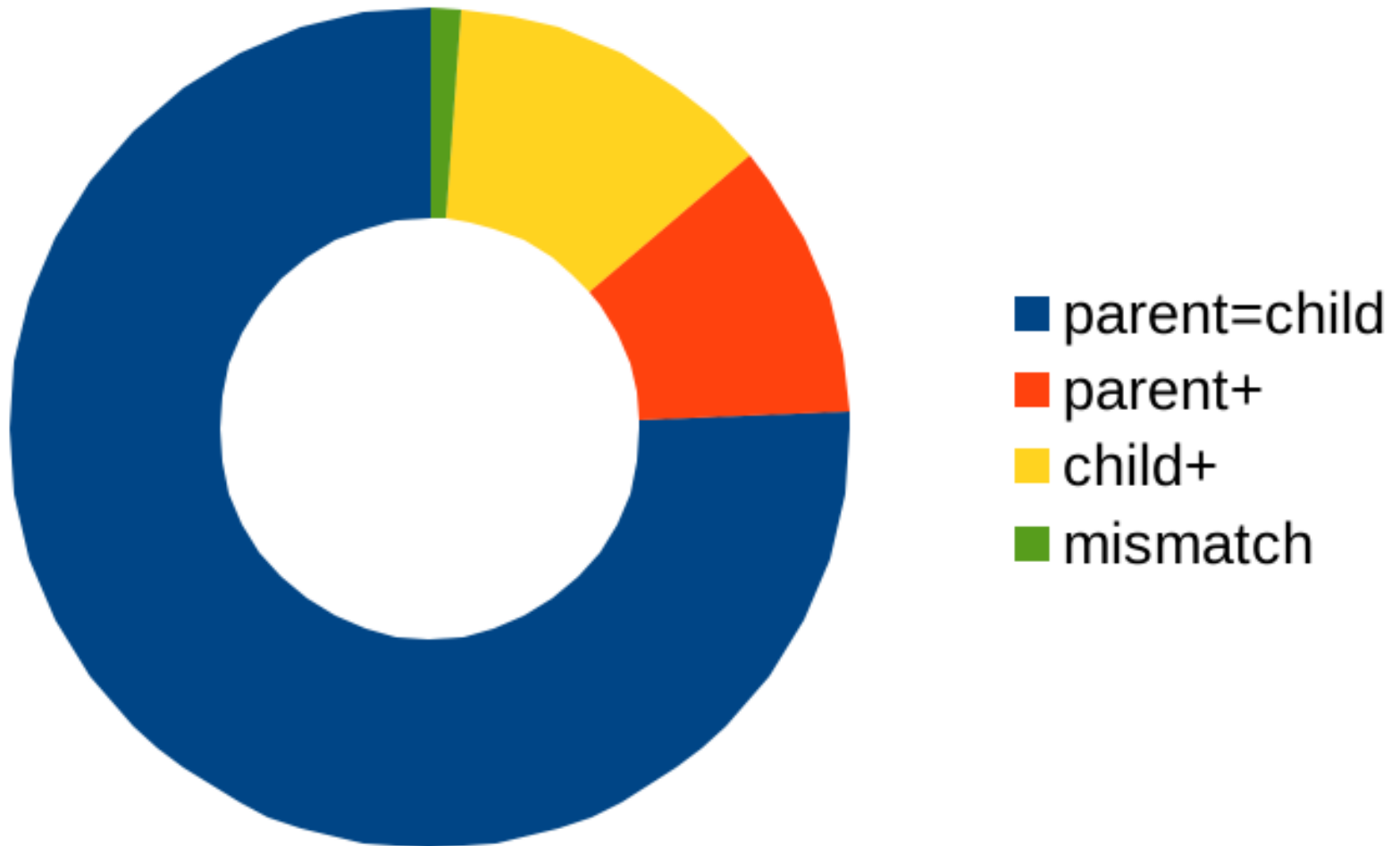
| | |
|------|-------------------------------------|
| name | IPv4 address |
| name | IPv4 address set |
| name | IPv6 address |
| name | IPv6 address set |
| name | IPv4 address + IPv6 address |
| name | IPv4 address set + IPv6 address |
| name | IPv4 address + IPv6 address set |
| name | IPv4 address set + IPv6 address set |

Evaluating Inconsistency

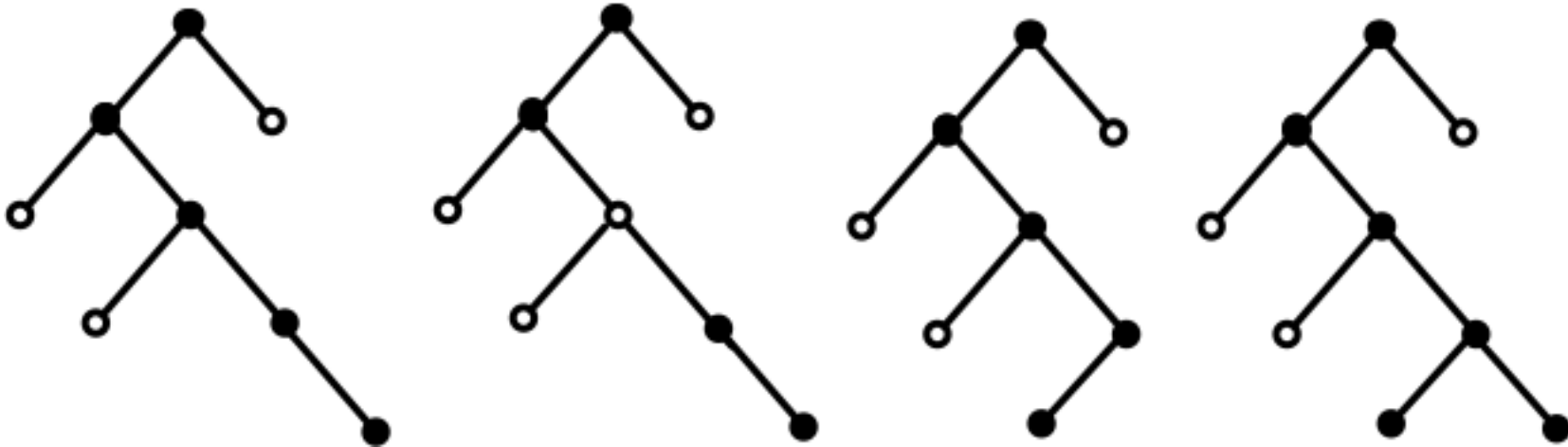
| | | |
|-------|---|------------|
| error | bad type (e.g. CNAME) | |
| error | bad rdata (e.g. Ipaddr for NS) | 29 (0.01) |
| error | TTL disagreement in NS RRset | 141 (0.06) |
| error | DNSSEC validation failure | |
| error | timeout/unreachable transient (e.g. down time) | |
| error | timeout/unreachable permanent (e.g. misconfiguration) | 1403 (6) |

| | | |
|----------------|-----------------------------------|------------|
| query_response | NOERROR | 21593 (90) |
| query_response | NXDomain | 23 (0.01) |
| query_response | REFUSED | 679 (3) |
| query_response | SERVFAIL / FORMERR / NOTIMP / ... | 142 (0.06) |
| query_response | referral after a referral | 77 (0.03) |
| query_response | aa==0 when aa==1 expected | 977 (4) |
| query_response | malicious or incorrect data | |

Parent/Child NS RRset Consistency



Namespace != Infrastructure Graph



Resolver (in)Stability

| | distribution | avoidance | recovery |
|------------|--------------|-----------|----------|
| BIND | proportional | no | < 1 sec |
| PowerDNS | spike dist. | no | 3 min |
| Unbound | uniform | yes | 15 min |
| DNSSCache | uniform | no | < 1 sec |
| WindowsDNS | uniform | yes | 1 sec |

- Source: Yu et al., Authority server selection in DNS cache resolvers, ACM SIGCOMM CCR 2012
- NOTE: negative caching => bursts of repeated failures

Discussion

- Inconsistency increases down the name space
- Inconsistency could exacerbate security threats
- Inconsistency may affect performance
- Inconsistency may lead to non-determinism

Questions

- Are some NS infrastructure graphs unknowable?
- Should consistency be encouraged? If so, how?
- There is no up/down sync, should there be?
- Should minimal-responses be preferred?
- Should repeated failures influence retry algorithms?
- Should NS RRs have had IPaddrs as RDATA?
- Is inconsistency worth studying further?